

# Linux (for PlayStation<sup>®</sup>2) VCL Preprocessor User's Manual

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Sony Computer Entertainment Inc.  
1-1, Akasaka 7-chome, Minato-ku  
Tokyo 107-0052, Japan

Sony Computer Entertainment America  
919 E. Hillsdale Blvd.  
Foster City, CA 94404, U.S.A.

Sony Computer Entertainment Europe  
30 Golden Square  
London W1F 9LD, U.K.

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## About This Manual

This manual provides a description of the various functionalities of the VCL Preprocessor (v1.3), for use with Linux (for PlayStation®2).

## Changes Since Last Release

None

## Typographic Conventions

Certain Typographic Conventions are used throughout this manual to clarify the meaning of the text:

Convention	Meaning
<code>courier</code>	Indicates literal program code.
<i>italic</i>	Indicates names of arguments and structure members (in structure/function definitions only).
<b>medium bold</b>	Indicates data types and structure/function names (in structure/function definitions only).
<a href="#">blue</a>	Indicates a hyperlink.

---

## Overview

### What is VCL?

The VCL preprocessor is an application that was developed to simplify some of the complex and tedious tasks associated with assembly-level programming of the VU processor. These tasks include:

- Dual pipeline processing
- Loop unrolling
- Register allocation
- Instruction scheduling

VCL outputs a standard VSM/DSM file (that can be compiled using ee-dvp-as). It is available for both the Linux and Win32 platforms.

### Merging of Upper and Lower Instructions in One Code Stream

VCL simplifies VU programming by merging upper and lower instructions. Pairing of instructions is no longer required.

### Syntax Simplification

In some cases, standard VU programming requires the programmer to specify which register will be used as a parameter in instructions such as MULQ. VCL performs the proper instruction assignment, and only requires the programmer to specify the more generic MUL. The rest is deduced from parameters attached to the instruction in question.

### Variable Naming and Registers Allocation

VCL provides for variable naming. Self-explanatory variable names like 'vertexptr' and 'vertexcolor' are permissible, instead of using standard register names such as 'v102' or 'vf06'.

### Instruction Scheduling

Assembly language programming for high-performance applications requires intimate knowledge of instruction timing (throughput and latency). Because of this, instruction scheduling is often very time consuming.

Data tracking is a tedious task because error messages regarding suspicious variable-related issues (use before set, for example) are not provided. High-performance code can appear to be confusing, especially if two blocks of code not otherwise related are reorganized into one block, for speed purposes.

While knowledge of instructions timing is still highly recommended, use of VCL will simplify instruction scheduling, keeping the code in a logical order, without compromising performance. VCL is aware of timing and dependencies, and in most cases will generate code that rivals hand-tuned code.

One of the most powerful features of VCL is related to loop unrolling. (See "Loop Unrolling" for more details.)

## Macro Usage

With C-preprocessor or GASP, macro usage has always been possible, even without VCL. Use is limited however, as instructions must be paired, and that this pairing may not be broken down. However, because VCL programming is single-streamed, it reinforces the power offered by macros.

## Syntax Simplification

VCL offers two syntax schemes, simply referred to as “old” and “new”. The “old” scheme is standard VU programming, where instructions must be specified using the full name, and where fields must be specified in instructions, as well as the registers to which they belong.

The “new” one attempts to simplify coding and code readability. To enable its use, the following must be added to the source file before any instructions:

```
.syntax new
```

An alternative way of enabling it is to specify “-n” as a command-line argument to VCL.

Note that merging of upper and lower instructions, as well as register naming, happens with either syntax. This section describes aspects of the “new” syntax.

### Merging of Upper and Lower Instructions

Because VCL manages instruction scheduling, it is no longer necessary to pair instructions. All instructions (from either the upper or lower pipeline) can now be ordered sequentially, in a single stream.

### Variable Naming and Registers Allocation

The use of named variables is permissible and encouraged. Besides making code more readable, named variables allow VCL to manage register allocation.

It is allowable to tie a specific register to a named variable and use register names directly. (See “Instruction Scheduling” for more details.) However, this limits the register allocation process.

The following are acceptable examples of variable naming usage:

```
IADDIU    inputptr, vi00, 32
LQ        vertex, 0(inputptr)
MAX       vector1111, vf00, vf00[w]
```

VCL checks to see if a variable is used before being set, and will output an error message if that is the case.

### Floating-Point and Integer Variable Naming

VCL tracks floating-point and integer variable names separately. Although not recommended, it is still possible to use the same variable name for both the floating point and the integer. Note that each will have to be initialized separately prior to use.

### Special Variable Name “i”

To ease the porting of C code to VCL, the variable name “i” is permitted for an integer register. Assuming such a named variable isn’t used with upper instructions, VCL will be able to discern between the hardware register I, and the variable. However, use of this variable name is discouraged, as it inhibits one of VCL’s primary functions, which is to make code more readable.

### Number Literals Peculiarities

In cases where a number literal must be specified, it is possible to specify instead a string, which will be assumed to be a defined value, and will therefore be ported as-is to the output file. Care must be taken, as VCL might not always be able to differentiate a typo from valid code. As an example, if you want to write:

## 4 Syntax Simplification

```
LOI                0x3F
```

but write instead:

```
LOI                x3F
```

VCL will accept it, only to have EE-DVP-AS reject it, with what could look like a cryptic message.

### Register Availability

It is necessary to let VCL know which registers are available for it to use. This is accomplished by using the following keywords:

```
.init_vi          VIxx <,VIxx>  
.init_vf          VFxx <,VFxx>  
.init_vi_all  
.init_vf_all
```

“.init\_vi” and “.init\_vf” are used to specify, respectively, which integer and floating-point registers are available. Specifying “vi00” or “vf00” is illegal, and will result in an error message. “.init\_vi\_all” and “.init\_vf\_all” may be alternatively used in lieu of specifying every single register by hand. Having both “.init\_vi” and “.init\_vi\_all”, for example, is illegal.

### Instruction Simplification

Many VU instructions stem from a single, more generic instruction. Such is the case for ADD, ADDi, ADDq, ADDbc, ADDA, ADDAi, ADDAq. In this case, the stem instruction would be ADD.

VCL accepts the replacement of the specific instructions by the stem instruction, as the specific instruction may be deduced by the parameters attached to the instruction. Therefore, a program could use:

```
ADD                vertexcolor, vertexcolor, q
```

And VCL would convert this (possibly) to:

```
ADDq.xyzw         vf03xyzw, vf03xyzw, q
```

The following table lists all instructions affected by the syntax simplification:

**Table 1**

Original Instruction	Simplified stem instruction
ADD	ADD
ADDi	
ADDq	
ADDbc	
ADDA	
ADDAi	
ADDAq	
ADDAbc	

Original Instruction	Simplified stem instruction
SUB	SUB
SUBi	
SUBq	
SUBbc	
SUBA	
SUBAi	
SUBAq	
SUBAbc	
MUL	
MULi	
MULq	
MULbc	
MULA	
MULAi	
MULAq	
MULAbc	
MADD	MADD
MADDi	
MADDq	
MADDbc	
MADDA	
MADDAi	
MADDAq	
MADDAbc	
MSUB	
MSUBi	
MSUBq	
MSUBbc	
MSUBA	
MSUBAi	
MSUBAq	
MSUBAbc	
MAX	MAX
MAXi	
MAXbc	
MINI	MINI
MINIi	
MINIbc	

## Floating-Point Register Fields Specification

VCL only requires that a floating-point register field be specified next to the instruction itself, as opposed to specifying it on the instruction and the register that it belongs to. Therefore, the following would be used:

```
ADD.xyz      newvertexposition, vertexposition, translation
```

Instead of:

```
ADD.xyz      vf04xyz, vf03xyz, vf02xyz
```

Specifying none is understood to be the same as specifying all (xyzw). Therefore, the following 2 cases would be equivalent:

```
MAX          color, color, vector0000
MAX.xyzw     color, color, vector0000
```

## Broadcast Instructions

When a field used as a broadcast is specified in the instruction (such as ADDbc or MADDbc), VCL requires the specified field to be next to the associated register. The following example is a typical vertex multiplication by a matrix.

```
MUL  acc,          matrix0, inputvertex[x]
MADD acc,          matrix1, inputvertex[y]
MADD acc,          matrix2, inputvertex[z]
MADD finalvertex, matrix3, inputvertex[w]
```

---

## Instruction Scheduling and Data Tracking

### Instruction Scheduling

Another useful feature of VCL is instruction rescheduling, which maximizes execution speed and code compactness. VCL tracks timing for each instruction (throughput and latency), and will try to reorder instructions to minimize stalls and maximize efficiency. In some rare cases (where possible), it will also try to move instructions from the upper to lower pipeline, and vice-versa.

VCL also tracks the I, ACC, Q and P registers, as well as the CLIP flags, and will generate proper delays between related instructions to insure valid code generation.

If, for some reason, you want to avoid instruction rescheduling past a certain point, simply insert the keyword:

```
--barrier
```

**Important:** Under certain circumstances, using the following line of code in a VSM/DSM file causes a problem:

```
NOP[E]          XGKICK          VIxx
```

To avoid the problem, placing these two instructions on two different lines:

```
NOP             XGKICK          VIxx
NOP[E]          NOP
```

To ensure such code isn't generated, follow an XGKICK instruction with "--barrier".

### Loop Unrolling

Instead of repeating the same code many times in a row, VCL allows loop unrolling by specifying the following, at the beginning of a loop:

```
--LoopCS      n,m
```

The loop ends when a branch instruction to the beginning of the loop is encountered. The following would constitute a valid loop:

```
LoopStart:
--LoopCS 3,3

LQI          inputvertex, 0(inputptr++)

MUL          acc,          matrix0, inputvertex[x]
MADD         acc,          matrix1, inputvertex[y]
MADD         acc,          matrix2, inputvertex[z]
MADD         finalvertex, matrix3, inputvertex[w]

SQI          finalvertex, 0(outputptr++)

IBNE        inputptr, endbuffer, LoopStart
```

At this time, loop unrolling is limited to simple loops without conditional branching. While VCL will not fail or give any error message if a branch is encountered inside the loop, the code generated will be less than optimal.

If the loop in question is performing clipping operations, it is possible to set the ADC bit without actually doing any conditional branch:

```
CLIPW.xyz    clipvtx, ClipData    ; Trigger clip calculations
FCAND       vi01, 0x3FFFF        ; Set if any of previous 3 vtx is clipped
IADDIU      adc_bit, vi01, 0x7FFF
ISW.w       adc_bit, outvl_xyzf2(output_buffer)    ; Set if clipped
```

Loop unrolling almost always results in a prologue followed by the main loop body, then concluded by an epilogue. The size of the prologue and epilogue depends partly on the parameters given to “--LoopCS”, but also depends on how well the code can be rescheduled.

The “--LoopCS” keyword takes two parameters: n (minimum number of loops) and m (slop count).

### n (Minimum Number of Loops)

As it unrolls the loop, VCL needs to know the loop’s minimum iteration count. This will allow it to potentially move instructions with side effects (stores, for example) higher in the execution pipeline, sometimes even before any conditional branch is executed. This translates into tighter (hence faster) code.

### m (Slop Count)

The slop count describes how many output iterations can be done without overwriting data past the end of an output array. A value of 1 would indicate that it is safe to execute an output instruction one iteration ahead of the current iteration, allowing for better instructions scheduling in some cases.

If the input is 30 vertices and “--LoopCS 3, 3” is specified, a 33-vertex output buffer will be required. If the input is less than 30 vertices, specify “--LoopCS 3, 0” instead. Note that the VCL may not process vertices ahead. In this case it will not matter if m is equal to 3, or even 500. It will behave the same as if m is 0.

**Tip:** The number of loop iterations is often based on a counter, which is decremented once per iteration. The loop is repeated until that counter reaches 0. A vertex counter is a good example of this. If the loop happens to be running through a given array, it is worth nothing that instead of using a counter, the address of the buffer ending may be calculated ahead of the loop. Then, instead of comparing the counter to 0, compare the current array pointer to the end pointer for equality. The end result is a saved instruction (the decrementing of the counter), tighter loop, and therefore faster code.

(Refer to Appendix B for more information.)

## Instructions Ordering

Care must be taken to place instructions in the order in which they are to take place. For example, standard VU programming would allow a DIV instruction right above an MULQ. The result of the previous DIV would be used, as it is known that the result of a DIV instruction isn’t available for seven cycles. With VCL, the MULQ must be placed before DIV. VCL will reposition them as it sees fit, to maximize execution speed and code compactness.

### Memory Aliasing and Instructions Reordering

VCL reorders instructions while preserving logical order. Instructions modifying a given variable will always appear in the same relative order to each other as they are in the input file. However, VCL has no explicit knowledge of memory, and more specifically, of memory aliasing.

Memory aliasing occurs when there are two different pointers potentially pointing to the same area in memory. In such a case, it is important that instructions reading and writing to both pointers’ memory be kept in the same relative order.

Steps must be taken to let VCL know that two pointers are potentially referring to the same memory. By simply appending a suffix to all instructions related to the aliased pointers, it will know to preserve the relative order. The following is an example for “ptr1” and “ptr2”:

```
SQ    var1, 0(ptr1):memgroup1
LQ    var2, 3(ptr2):memgroup1
```

Here, “memgroup1” could be any valid string. It ensures that the store (SQ) always takes precedence to the load (LQ).

Within one program, more than one group may be used.

### Peculiarities with XGKick

Moving XGKick instructions could potentially result in hazardous code. For example, the XGKick could be moved ahead of stores to the buffer to be XGKicked. To avoid any such cases, XGKick will never be moved before or after a store instruction.

### Branch Delay Slots

VCL handles branch delay slots independently. Placing any instruction immediately after a branch will cause the instruction in question to be executed only if the branch is not taken (in the case of a non-returning branch) or when the program counter comes back from a sub-routine (in the case of a function call such as BAL).

### Code Removal

VCL will recognize if a code block isn't reachable and will remove it. This saves VU micro memory in the process.

Empty lines (NOP NOP) are also removed, and will be replaced if necessary by either a WAITQ or WAITP instruction. Any stall introduction by instruction removal is noted by VCL as a comment in the output file.

### Floating-Point Field Pruning

VCL keeps track of which fields of a floating-point register are used, and will prune any field that does not need to be used. This step allows VCL to more effectively schedule instructions under some circumstances. The following is an example of pruning:

```
LQ    color, 2(inputptr)
ADD   color, color, ambientcolor
MINI  color, color, vector1111
SQ.xyz color, 2(outputptr)
```

Assuming that “color” isn't used later in the code, VCL will prune the w field in the three first instructions. The code effectively becomes the same as:

```
LQ.xyz color, 2(inputptr)
ADD.xyz color, color, ambientcolor
MINI.xyz color, color, vector1111
SQ.xyz color, 2(outputptr)
```

### E, D, and T Bits

Putting an [E], [D] and/or [T] bit on an instruction is valid with VCL. However, as it is rescheduled, the bit will move around with the instruction it is attached to. Also, if the instruction it is attached to is duplicated (in such cases as loop unrolling), the bit will be duplicated as well.

Instead of using the [E] bit, it is recommended to use the VCL keyword:

```
--cont
```

“--cont”, which stands for “continue”, lets VCL know the program will restart from this point. Effectively, VCL inserts a “NOP[E] NOP” line. The main advantage of using this keyword over inserting an explicit [E] bit is that it removes any danger of VCL moving the [E] bit somewhere unexpected. Another solution to this problem is to attach the bit to a label such as:

```
Label:[D]
```

Also note, as with “--barrier”, instructions are not rescheduled beyond a “--cont”.

## Load and Store Offsets

When a loop containing load and store instructions LQ and SQ is unrolled, VCL will most likely alter the offset parameter to at least some of the instructions. This is a result of moving the instruction past the increment of the pointer the offset relates to. The following is an example:

```
LQ          vertex, 0(inputptr)
LQ          stq,   1(inputptr)
LQ          rgb,   2(inputptr)
IADDIU     inputptr, inputptr, 3
...
```

A possible (and simplistic) unrolling of such an instruction sequence would be:

```
LQ          vertex1, 0(inputptr)
LQ          stq1,   1(inputptr)
LQ          rgb1,   2(inputptr)
IADDIU     inputptr, inputptr, 3
...
LQ          vertex2, 0(inputptr)
LQ          stq2,   1(inputptr)
LQ          rgb2,   2(inputptr)
IADDIU     inputptr, inputptr, 3
...
LQ          vertex3, 0(inputptr)
LQ          stq3,   1(inputptr)
LQ          rgb3,   2(inputptr)
IADDIU     inputptr, inputptr, 3
...
```

However, VCL might instead unroll a sequence as follows, which produces the same result, but under certain conditions provides better instruction rescheduling:

```
LQ          vertex1, 0(inputptr)
LQ          stq1,   1(inputptr)
IADDIU     inputptr, inputptr, 3
LQ          rgb1,   -1(inputptr)
...
LQ          vertex2, 0(inputptr)
LQ          rgb2,   2(inputptr)
IADDIU     inputptr, inputptr, 3
...
LQ          stq2,   -2(inputptr)
LQ          vertex3, 0(inputptr)
LQ          stq3,   1(inputptr)
LQ          rgb3,   2(inputptr)
IADDIU     inputptr, inputptr, 3
...
```

Note that using LQI, LQD, SQI, and SQD will prevent the possibility of such optimizations. However, their use has proven to be more effective under certain conditions such as very tight loops.

## Data Tracking

VCL tracks data usage for many reasons. If VCL determines that a variable is set but not used afterwards, it will remove the setting instructions. If these settings need to remain in the code for the purpose of eventually passing the variable as a parameter out of the code block, use the following keyword:

```
out_vi (VIxx) intvarname
out_vf (VFxx) floatvarname
```

“VIxx” and “VFxx” are an integer and a float register, respectively. The names “intvarname” and “floatvarname” correspond to these registers. VCL will make sure the specified register will contain the named variable value. These keywords must appear between the two keywords “--exit” and “--endexit”, or between “--exitm” and “--endexit”. (See “-exit / -endexit” for more details.)

Note that similar instructions exist for ACC (out\_hw\_acc acc), I (out\_hw\_i i), P (out\_hw\_p p), Q (out\_hw\_q q) registers, and the CLIP flags (out\_hw\_clip clip).

## Set Before Use

If a variable is found to be used before it is even initialized, VCL will output an error message. If it is necessary to pass in a variable as a parameter from a different block, or from standard VSM/DSM code, use the following keywords:

```
in_vi (VIxx) intvarname
in_vf (VFxx) floatvarname
```

These keywords must appear between the two keywords “--enter” and “--endenter”. (See “-enter/-endenter” for more details.)

Specifying “in\_vi” and “in\_vf” does not automatically mean the link between the register and the variable name will remain for the rest of the program, even if the same link is specified using “out\_vi” or “out\_vf”. VCL reserves the right to break that link at any point.

Note that similar instructions exist for ACC (in\_hw\_acc acc), I (in\_hw\_i i), P (in\_hw\_p p) and Q (in\_hw\_q q) registers, as well as the FPU results flags (in\_hw\_status status) and the CLIP flags (in\_hw\_clip clip).

An alternative to using “in\_hw\_clip clip” is to have the following instruction before any clipping-related instruction:

```
FCSET 0
```

This will let VCL know all clipping flags are initialized.

VCL does not keep track of all flags. It is aware of MAC flags, but not about sticky bits.

To initialize a variable to 0, it is illegal to use the following:

```
SUB varname, varname, varname
```

Instead, use one of the two following methods:

```
SUB varname, vf00, vf00
MFIR varname, vi00; Flags are left untouched!
```

---

## Branching

### Labels

With VCL, labels are defined in the same way as under VSM/DSM. However, for a label to remain in the output file, it must be positioned between the “--enter”/“--endenter” and “--exit”/“--endexit” keywords (or “--exitm”/“--endexit”). If “--exit”/“--endexit” is omitted, the label can be positioned after “--enter”/“--endenter”.

### Calls to Functions

VCL currently does not support far function calls (calls to external functions). Therefore, any functions called must be included in the same file as the caller, either directly or via a header file (for `c-preprocessor`) or include file (for `gasp`).

Within VCL, a function is treated like any branch. There is no special method of passing parameters in and out. The following code is an example:

```

        IADDIU      vertexptr, vi00, 64
        IADDIU      endptr, vertexptr, 20

        BAL        retaddress, TransformVerticesInPlace

        ...

TransformVerticesInPlace:
        LQ          vertex, 0(vertexptr)

        ;*** Matrix used without being initialized! ***
        MatrixMultiplyVertex vertex, matrix, vertex

        SQI          vertex, 0(vertexptr++)
        IBNE        vertexptr, endptr

        JR          retaddress

```

In this case (for register scope), the function “`TransformVerticesInPlace`” behaves the same as if it were in-lined. Assuming the matrix hasn’t been initialized before the function call, an error would be given to this effect.

### Functions Calling Sub-Functions

It is valid for a function to call a sub-function. However, be aware that returning addresses conflicts will result in an infinite loop.

## Recursive Functions

VCL does not support recursive functions natively. However, it is possible for a program to maintain its own stack, push the return address register before calling the function recursively, pop the address on return, and eventually return, as is shown in the example below:

```
RecursiveFunction:
...

ISUBIU      counter, counter, 1
IBEQ       counter, vi00, RecFcnEnd

ISUBIU      stackptr, stackptr, 1
ISWR.x     retaddress, (stackptr):VCLSML_STACK

BAL        retaddress, RecursiveFunction):VCLSML_STACK

ILWR.x     retaddress, (stackptr):VCLSML_STACK
IADDIU     stackptr, stackptr, 1

RecFcnEnd:
JR         retaddress
```

## Jump Tables

VCL currently does not support jump tables. Such feature must therefore be partly implemented in standard VSM/DSM code. However, each function of the jump table may be coded separately using VCL.

---

## Integration of VSM Code Within VCL

Sometimes it may be necessary to incorporate traditional VSM/DSM code within VCL code. This section introduces two methods to do so, and explains how they are used.

### **.vsm / .endvsm and .raw / .endraw**

This is used for inserting pre-formatted, pre-ordered VSM code. Instructions therefore have to be organized in the original 2 stream (upper and lower). However, variables may still be named, and use the same ones as non-VSM code. Checks for data use before set will still be performed.

“.raw” / “.endraw” is simply an alternative spelling for “.vsm” / “.endvsm”.

Note that code within a VSM block isn't rescheduled. The block acts the same way as “--barrier”, in that instruction rescheduling is not performed across the block. You must consider whether or not the use of such a block contributes to better overall performance.

Also note that because such block is meant for code only, it must appear between “--enter” / “--endenter” and “--exit” / “--endexit” (or “--exitm” / “--endexit”). It must not be thought of as a black box that is ported as-is to the output file. For the same reason, specifying keywords like “.equ” inside such a block is invalid.

The following is an example of a .vsm / .endvsm block:

```
IADDUI      matrixptr, vi00, 0
MatrixLoad  mat1, 0, (matrixptr)
MatrixLoad  mat2, 4, (matrixptr)

; Swap matrices
.vsm
max mat1[0], mat2[0], mat2[0]   move mat2[0], mat1[0]
max mat1[1], mat2[1], mat2[1]   move mat2[1], mat1[1]
max mat1[2], mat2[2], mat2[2]   move mat2[2], mat1[2]
max mat1[3], mat2[3], mat2[3]   move mat2[3], mat1[3]
.endvsm

MatrixSave  mat1, 0, (matrixptr)
MatrixSave  mat2, 4, (matrixptr)
```

### **.rawloop / .endrawloop**

These directives are used to enclose a pre-formatted VSM code block to be unrolled. More control is permitted as far as prologue and epilogue creation goes. There following is an example:

```

.rawloop
loop:
--LoopCS 10,10

5..ftoi4.xyz  ixyz, sxyz          1..lqi   vrt, (ptr++)
4..mul.xyz    sxyz, nxfrm, q      nop
               nop              nop
1..mul        acc, m[3], vf00[w]  2..move  nxfrm, xfrm
1..madd       acc, m[0], vrt[x]   2..div   q, vf00[w], xfrm[w]
1..madd       acc, m[1], vrt[y]   1..ibne  ptr, end_ptr, loop
1..madd       xfrm, m[2], vrt[z]  5..sqi  ixyz, (optr++)

.endrawloop

```

The “--LoopCS” keyword is used the same way as described in “Loop Unrolling”. While building the prologue, the instructions with the lowest numbers are introduced first, and follow this order: 1 1-2 1-2-3 1-2-3-4 ... etc. Then the loop itself and the epilogue are created in accordance with the prologue.

(Refer to Appendix B for more information.)

---

## Macros and Other Preprocessor Usages

Using macros with VU code has always been possible, using tools such as the C preprocessor and `gasp`. The use of macros is also possible with VCL. In fact, VCL offers even more opportunities to use these tools, as the input is single-streamed, as opposed to the double-streamed VSM/DSM programming style.

### Using the C Preprocessor

VCL can automatically pipe the code through the C preprocessor by giving it the “-G” (uppercase G) command-line parameter.

Using it will allow the use of all preprocessor directives:

```
#if, #ifdef, #ifndef
#elif, #else
#endif
#include
#define
#undef
#line
#error
```

### Using Macros with the C Preprocessor

Use of macros with the C preprocessor is identical to that of C/C++. However, using macros with a label inside proves to be an interesting task. The only way around doing this is to include an extra parameter to the macro, which will then be used internally to generate a unique label, using the `##` token concatenation directive:

```
#define mymacro(param1,param2,uniqueid)      \
{                                             \
    IADDIU    Counter, vi00, 5              \
                                             \
    mymacrolabel ## uniqueid:              \
    ISUBIU    Counter, Counter, 1          \
    IBNE      Counter, vi00, mymacrolabel ## uniqueid \
}                                             \
```

Each call to “`mymacro`” within a given file (including all the included header files) must pass a unique identifier for “`uniqueid`”.

### Using `gasp`

Just as with the C preprocessor, VCL can automatically pipe the code through `gasp`, by giving it the “-g” (lowercase g) command-line parameter.

Using `gasp` will allow the use, among others, of the following directives:

```
.assign
.include
.macro
.endm
.end
```

Refer to the `gasp` documentation for more information on the subject. Online documentation may be found at the following URLs:

[http://www.objsw.com/docs/gasp\\_toc.html](http://www.objsw.com/docs/gasp_toc.html)

[http://sunsite.utk.edu/gnu/binutils/gasp\\_toc.html](http://sunsite.utk.edu/gnu/binutils/gasp_toc.html)

[http://case.ispras.ru/PublicScripts/cgi-bin/lib.cgi/gnu/gasp\\_toc.html](http://case.ispras.ru/PublicScripts/cgi-bin/lib.cgi/gnu/gasp_toc.html)

## Using Macros with `gasp`

As shown in the following example, using macros with `gasp` is simple:

```
.macro mymacro param1,param2
IADDIU    Counter, vi00, 5
mymacrolabel\@:
ISUBIU    Counter, Counter, 1
IBNE      Counter, vi00, mymacrolabel\@
.endm
```

`gasp` will replace the “\@” by a number, which is incremented with each instance of a macro, so the same macro may be used many times in the same source file.

## Issues with `gasp`

With some versions of `gasp`, including `gasp`, number literals aren’t translated properly to the output file. Examples of failed conversions are:

```
0x123      (converted to)  0
0.000123   (converted to)  0.1
```

The way to fix this, for now, is to use the following syntax:

```
0x123      (switch to)    H'123
0.000123   (switch to)    1.23E-4 or (0.0001)
```

Also, any line without a space or tab in front of the first non-white character will be converted as though the first word was a label. The newly created label will be suffixed by a semi-column.

## Examples of Preprocessor Usage

Refer to Appendix A for examples.

---

## Command-Line Parameters

VCL is a command-line-based preprocessor. Various parameters may be passed to it, all of which are described in this chapter.

### Command-Line Syntax

VCL must be called with parameters, following the syntax:

```

vcl [-cCdefgGhKlMmNnPSZ] [-I<includefilepath>] [-t<seconds>]
    [-o<outputfilename>] [-u<string>] <inputfilename>

```

#### **-c**

Emit nearly original source code as comments.

#### **-C**

Disable the code reduction pass.

#### **-d**

Dumb code is generated. For example, rescheduling of instructions isn't performed.

#### **-e**

Disable the generation of [E] bits at the end of the code. Alternatively, the use of `-exitm` without an argument may be used.

#### **-f**

Disable the generation of alignment directives (`.align n`).

#### **-g**

Run `gasp` on the input before any VCL-specific task is done. `gasp` is called with the following parameter string: `"-p -s -c ';"`. `-l` is also passed if specified.

#### **-G**

Run the C preprocessor on the input before any VCL-specific task is done.

#### **-h**

Print out the command-line help.

#### **-I<includefilepath>**

To be used with `"-g"`; tells `gasp` where to find include files.

**<inputfilename>**

Specify the name of the VCL source file. If it is not specified, VCL will read from the standard input.

**-K**

The temporary files created by the pre-processors are not deleted. The file locations being OS-dependant, refer to the VCL output to find them.

**-L**

Globally disable loop code generation.

**-m**

Generate “.mpg” and DMA tags automatically. This may be used as an alternative to “.mpg” within the VCL source file.

**-M**

VCL retains the relative order of load and store instructions (known as the timid memory access mode).

**Note:** -M will be deprecated in future versions.

**-n**

Enable the new syntax. This may be used as an alternative to “.syntax new” within the VCL source file. (See “Syntax Simplification” for more details.)

**-o<outputfilename>**

Output file name. Specifying the same name as the source is invalid. If not specified, the result is outputted to the standard output.

**-P**

Disable the removal of unused instructions and pruned fields pass. (See “Code Removal” for more details.)

**Note:** -P will be deprecated in future versions.

**-S**

The content of loops starting with “--LoopCS” will be reorganized to stagger the read and write instructions, and to facilitate memory access by the VIF and GIF.

**-t<seconds>**

Specify the optimizer timeout. <second> must be 1 or higher. Default is 4. Note that this has the potential side-effect of generating different code on different computers, as the processor speed is not taken into account.

**-u<string>**

<string> is used as a unique string for label generation, instead of the file name. Useful if the filename is especially long.

## **-Z**

Disable the immediate field fix up pass.

**Note:** -Z will be deprecated in future versions.

---

## Keywords

### **.global symbolname**

The directive “.global”, along with “symbolname”, is ported as-is to the output file. “symbolname” is therefore only assumed to exist and be valid.

### **.init\_vi Vlxx <, Vlxx ...>**

Inform VCL that the specified integer registers are available for use. Specifying VI00 is illegal, as it is always considered available. (Specifying “.init\_vi” and “.init\_vi\_all” in the same file is illegal. Note that VCL might fail to process code if not given enough registers.

### **.init\_vf VFxx <, VFxx ...>**

Inform VCL that the specified float registers are available for use. Specifying VF00 is illegal, as it is always considered available (Compare to 8.4.). Specifying “.init\_vf” and “.init\_vf\_all” in the same file is illegal. Note that VCL might fail to process code if not given enough registers.

### **.init\_vi\_all**

Inform VCL that all integer registers are available for use. Specifying “.init\_vi” and “.init\_vi\_all” in the same file is illegal.

### **.init\_vf\_all**

Inform VCL that all float registers are available for use. Specifying “.init\_vf” and “.init\_vf\_all” in the same file is illegal.

### **.mpg vucodeoffset**

Add “ret” DMA tags around the code generated by VCL, for better integration of VCL code with original VSM/DSM code. If “.name” is also specified, two labels will be added, following the syntax “(progrname)\_DmaTag” and “(progrname)\_DmaEnd”. “vucodeoffset” is assumed to be a valid address.

### **.name progrname**

Add two labels, one before the code generated by VCL and the other after, following the syntax “(progrname)\_CodeStart” and “(progrname)\_CodeEnd”. For better integration of VCL code with original VSM/DSM code. The labels created are also made available globally, via the directive “.global”.

### **.raw / .endraw**

Enclose pre-formatted, original VSM-style code. Same as “vsm / .endvsm”. (See “.vsm / .endvsm and .raw / .endraw” for more details.)

### **.rawloop / .endrawloop**

Enclose a pre-formatted, original VSM-style code loop, to be unrolled. (See “.rawloop / .endrawloop” for more details.)

**.syntax old | new**

If “old” is specified, the syntax is the same as original VSM/DSM code. “new” specifies the new and simplified syntax. (See “Syntax Simplification” for more details.)

**.vsm / .endvsm**

Enclose pre-formatted, original VSM-style code. Same as “.raw” / “.endraw”. (See “.vsm / .endvsm and .raw / .endraw” for more details.)

**--barrier**

Prevent the rescheduling of instruction to go across this line. (See “Instruction Scheduling” for more details.)

**--cont**

Mark a point where a program temporarily stops, and may be restarted from, via a MSCNT. A [E] flag is inserted at this point. (See “E, D and T Bits” for more details.)

**--enter / --enderter**

Specify an entry point to VCL code. Any file must have at least one entry point, but may have more than one.

**in\_vi (Vlxx) varname**

Must be specified between “--enter” and “--enderter”. Bind a specific integer to a specific variable name at entry. The register is considered pre-initialized, presumably by standard VSM/DSM code. Such binding is not guaranteed to persist for the duration of the code block. (See “Set Before Use” for more details.)

**in\_vf (VFxx) varname**

Must be specified between “--enter” and “--enderter”. Bind a specific float to a specific variable name at entry. The register is considered pre-initialized, presumably by standard VSM/DSM code. Such binding is not guaranteed to persist for the duration of the code block. (See “Set Before Use” for more details.)

**in\_hw\_acc acc / in\_hw\_clip clip / in\_hw\_i i / in\_hw\_p p / in\_hw\_q q / in\_hw\_status status**

Must be specified between “--enter” and “--enderter”. The specified register is considered pre-initialized, presumably by standard VSM/DSM code. (See “Set Before Use” for more details.)

**--exit / --endexit**

Specify an exit point to VCL code. Its use is mandatory only if outputting parameters is necessary.

**--exitm macroname / --endexit**

Specify an exit point to VCL code. Its use is mandatory only if outputting parameters is necessary. Unlike “--exit”, “--exitm” lets you specify a macro name, which will be sent as-is to the output file, therefore permitting custom ending code. (See “bla bla bla” for more details.)

**out\_vi (Vlxx) varname**

Must be specified between “--exit” and “--endexit” or between “--exitm” and “--endexit”. Bind a specific integer to a specific variable name at exit. The variable may then be passed to other VCL blocks, or to original VSM/DSM code. Such binding is not guaranteed to persist for the duration of the code block. (See “Set Before Use” for more details.)

**out\_vf (VFxx) varname**

Must be specified between “--exit” and “--endexit” or between “--exitm” and “--endexit”. Bind a specific float to a specific variable name at exit. The variable may then be passed to other VCL blocks, or to original VSM/DSM code. Such binding is not guaranteed to persist for the duration of the code block. (See “Set Before Use” for more details.)

**out\_hw\_acc acc / out\_hw\_clip clip / out\_hw\_i i / out\_hw\_p p / out\_hw\_q q**

Must be specified between “--exit” and “--endexit”, or between “--exitm” and “--endexit”. The specified register may then be passed to other VCL blocks, or to original VSM/DSM code. (See “Data Tracking” for more details.)

**--LoopCS n,m**

Mark a portion of code as being a loop, and instruct VCL to unroll it. “n” is the minimum iteration of the loop, and “m” (slop count) is the amount of output iterations that can be done without overwriting data past the end of an output array. (See “Loop Unrolling” for more details.)



## Appendix A: Macro Examples

All macros in this appendix may be found in the file VCL\_SML.i, included with the VCL distribution. All can be used as-is with `gasp` (via the “-g” command-line parameter), but all can easily be converted to be used by the C preprocessor.

```

;///-----
;/// MatrixLoad - Load "matrix" from VU mem location "vumemlocation" +
;/// "offset"
;///-----
    .macro    MatrixLoad    matrix,offset,vumemlocation
    lq        \matrix[0], \offset+0(\vumemlocation)
    lq        \matrix[1], \offset+1(\vumemlocation)
    lq        \matrix[2], \offset+2(\vumemlocation)
    lq        \matrix[3], \offset+3(\vumemlocation)
    .endm

;///-----
;/// MatrixSave - Save "matrix" to VU mem location "vumemlocation" +
;/// "offset"
;///-----
    .macro    MatrixSave    matrix,offset,vumemlocation
    sq        \matrix[0], \offset+0(\vumemlocation)
    sq        \matrix[1], \offset+1(\vumemlocation)
    sq        \matrix[2], \offset+2(\vumemlocation)
    sq        \matrix[3], \offset+3(\vumemlocation)
    .endm

;///-----
;/// MatrixIdentity - Set "matrix" to be an identity matrix
;/// Thanks to Colin Hugues (SCEE) for that one
;///-----
    .macro    MatrixIdentity    matrix
    add.x     \matrix[0], vf00, vf00[w]
    mfir.yzw  \matrix[0], vi00

    mfir.xzw  \matrix[1], vi00
    add.y     \matrix[1], vf00, vf00[w]

    mr32     \matrix[2], vf00

    max      \matrix[3], vf00, vf00
    .endm

;///-----
;/// MatrixCopy - Copy "matrixsrc" to "matrixdest"
;/// Thanks to Colin Hugues (SCEE) for that one
;///-----
    .macro    MatrixCopy    matrixdest,matrixsrc
    max      \matrixdest[0], \matrixsrc[0], \matrixsrc[0]
    move     \matrixdest[1], \matrixsrc[1]
    max      \matrixdest[2], \matrixsrc[2], \matrixsrc[2]
    move     \matrixdest[3], \matrixsrc[3]
    .endm

```

```

;/-------
; // MatrixSwap - Swap the content of "matrix1" and "matrix2"
; // The implementation seems lame, but VCL will convert moves to maxes
; // if it sees fit
;/-------
    .macro    MatrixSwap matrix1,matrix2
    move      vclsmftemp, \matrix1[0]
    move      \matrix1[0], \matrix2[0]
    move      \matrix2[0], vclsmftemp

    move      vclsmftemp, \matrix1[1]
    move      \matrix1[1], \matrix2[1]
    move      \matrix2[1], vclsmftemp

    move      vclsmftemp, \matrix1[2]
    move      \matrix1[2], \matrix2[2]
    move      \matrix2[2], vclsmftemp

    move      vclsmftemp, \matrix1[3]
    move      \matrix1[3], \matrix2[3]
    move      \matrix2[3], vclsmftemp
    .endm

;/-------
; // MatrixTranspose - Transpose "matrixsrc" to "matresult". It is safe
; // for "matrixsrc" and "matresult" to be the same.
; // Thanks to Colin Hugues (SCEE) for that one
;/-------
    .macro    MatrixTranspose matresult,matrixsrc
    mr32.y    vclsmftemp, \matrixsrc[1]
    add.z     \matresult[1], vf00, \matrixsrc[2][y]
    move.y    \matresult[2], vclsmftemp
    mr32.y    vclsmftemp, \matrixsrc[0]
    add.z     \matresult[0], vf00, \matrixsrc[2][x]
    mr32.z    vclsmftemp, \matrixsrc[1]
    mul.w     \matresult[1], vf00, \matrixsrc[3][y]
    mr32.x    vclsmftemp, \matrixsrc[0]
    add.y     \matresult[0], vf00, \matrixsrc[1][x]
    move.x    \matresult[1], vclsmftemp
    mul.w     vclsmftemp, vf00, \matrixsrc[3][z]
    mr32.z    \matresult[3], \matrixsrc[2]
    move.w    \matresult[2], vclsmftemp
    mr32.w    vclsmftemp, \matrixsrc[3]
    add.x     \matresult[3], vf00, \matrixsrc[0][w]
    move.w    \matresult[0], vclsmftemp
    mr32.y    \matresult[3], vclsmftemp
    add.x     \matresult[2], vf00, vclsmftemp[y]

    move.x    \matresult[0], \matrixsrc[0]           ;// These 4
instructions will be
    move.y    \matresult[1], \matrixsrc[1]           ;// removed if
"matrixsrc" and
    move.z    \matresult[2], \matrixsrc[2]           ;// "matresult" are
the same
    move.w    \matresult[3], \matrixsrc[3]           ;//
    .endm

```

```

;///-----
;/// MatrixMultiply - Multiply 2 matrices, "matleft" and "matright", and
;/// output the result in "matresult". Dont forget matrix multipli-
;/// cations arent commutative, i.e. left X right wont give you the
;/// same result as right X left.
;///
;/// Note: ACC register is modified
;///-----
    .macro    MatrixMultiply    matresult,matleft,matright
    mul      acc,                \matright[0], \matleft[0][x]
    madd     acc,                \matright[1], \matleft[0][y]
    madd     acc,                \matright[2], \matleft[0][z]
    madd     \matresult[0], \matright[3], \matleft[0][w]

    mul      acc,                \matright[0], \matleft[1][x]
    madd     acc,                \matright[1], \matleft[1][y]
    madd     acc,                \matright[2], \matleft[1][z]
    madd     \matresult[1], \matright[3], \matleft[1][w]

    mul      acc,                \matright[0], \matleft[2][x]
    madd     acc,                \matright[1], \matleft[2][y]
    madd     acc,                \matright[2], \matleft[2][z]
    madd     \matresult[2], \matright[3], \matleft[2][w]

    mul      acc,                \matright[0], \matleft[3][x]
    madd     acc,                \matright[1], \matleft[3][y]
    madd     acc,                \matright[2], \matleft[3][z]
    madd     \matresult[3], \matright[3], \matleft[3][w]
    .endm

;///-----
;/// LocalizeLightMatrix - Transform the light matrix "lightmatrix" into
;/// local space, as described by "matrix", and output the result in
;/// "locallightmatrix"
;///
;/// Note: ACC register is modified
;///-----
    .macro    LocalizeLightMatrix    locallightmatrix,matrix,lightmatrix
    mul      acc,                \lightmatrix[0], \matrix[0][x]
    madd     acc,                \lightmatrix[1], \matrix[0][y]
    madd     acc,                \lightmatrix[2], \matrix[0][z]
    madd     \locallightmatrix[0], \lightmatrix[3], \matrix[0][w]

    mul      acc,                \lightmatrix[0], \matrix[1][x]
    madd     acc,                \lightmatrix[1], \matrix[1][y]
    madd     acc,                \lightmatrix[2], \matrix[1][z]
    madd     \locallightmatrix[1], \lightmatrix[3], \matrix[1][w]

    mul      acc,                \lightmatrix[0], \matrix[2][x]
    madd     acc,                \lightmatrix[1], \matrix[2][y]
    madd     acc,                \lightmatrix[2], \matrix[2][z]
    madd     \locallightmatrix[2], \lightmatrix[3], \matrix[2][w]

    move     \locallightmatrix[3], \lightmatrix[3]
    .endm

```

```

;///-----
;/// MatrixMultiplyVertex - Multiply "matrix" by "vertex", and output
;/// the result in "vertexresult"
;///
;/// Note: Apply rotation, scale and translation
;/// Note: ACC register is modified
;///-----
    .macro    MatrixMultiplyVertex vertexresult,matrix,vertex
    mul      acc,          \matrix[0], \vertex[x]
    madd     acc,          \matrix[1], \vertex[y]
    madd     acc,          \matrix[2], \vertex[z]
    madd     \vertexresult, \matrix[3], \vertex[w]
    .endm

;///-----
;/// MatrixMultiplyVertex - Multiply "matrix" by "vertex", and output
;/// the result in "vertexresult"
;///
;/// Note: Apply rotation, scale and translation
;/// Note: ACC register is modified
;///-----
    .macro    MatrixMultiplyVertexXYZ1 vertexresult,matrix,vertex
    mul      acc,          \matrix[0], \vertex[x]
    madd     acc,          \matrix[1], \vertex[y]
    madd     acc,          \matrix[2], \vertex[z]
    madd     \vertexresult, \matrix[3], vf00[w]
    .endm

;///-----
;/// MatrixMultiplyVector - Multiply "matrix" by "vector", and output
;/// the result in "vectorresult"
;///
;/// Note: Apply rotation and scale, but no translation
;/// Note: ACC register is modified
;///-----
    .macro    MatrixMultiplyVector vectorresult,matrix,vector
    mul      acc,          \matrix[0], \vector[x]
    madd     acc,          \matrix[1], \vector[y]
    madd     \vectorresult, \matrix[2], \vector[z]
    .endm

;///-----
;/// VectorLoad - Load "vector" from VU mem location "vumemlocation" +
;/// "offset"
;///-----
    .macro    VectorLoad vector,offset,vumemlocation
    lq      \vector, \offset(\vumemlocation)
    .endm

;///-----
;/// VectorSave - Save "vector" to VU mem location "vumemlocation" +
;/// "offset"
;///-----
    .macro    VectorSave vector,offset,vumemlocation
    sq      \vector, \offset(\vumemlocation)
    .endm

```

```

;///-----
;/// VectorAdd - Add 2 vectors, "vector1" and "vector2" and output the
;/// result in "vectorresult"
;///-----
        .macro    VectorAdd    vectorresult,vector1,vector2
        add            \vectorresult, \vector1, \vector2
        .endm

;///-----
;/// VectorSub - Subtract "vector2" from "vector1", and output the
;/// result in "vectorresult"
;///-----
        .macro    VectorSub    vectorresult,vector1,vector2
        sub            \vectorresult, \vector1, \vector2
        .endm

;///-----
;/// VertexLoad - Load "vertex" from VU mem location "vumemlocation" +
;/// "offset"
;///-----
        .macro    VertexLoad    vertex,offset,vumemlocation
        lq            \vertex, \offset(\vumemlocation)
        .endm

;///-----
;/// VertexSave - Save "vertex" to VU mem location "vumemlocation" +
;/// "offset"
;///-----
        .macro    VertexSave    vertex,offset,vumemlocation
        sq            \vertex, \offset(\vumemlocation)
        .endm

;///-----
;/// VertexPersCorr - Apply perspective correction onto "vertex" and
;/// output the result in "vertexoutput"
;///
;/// Note: Q register is modified
;///-----
        .macro    VertexPersCorr    vertexoutput,vertex
        div            q, vf00[w], \vertex[w]
        mul            \vertexoutput, \vertex, q
        .endm

;///-----
;/// VertexPersCorrST - Apply perspective correction onto "vertex" and
;/// "st", and output the result in "vertexoutput" and "stoutput"
;///
;/// Note: Q register is modified
;///-----
        .macro    VertexPersCorrST    vertexoutput,stoutput,vertex,st
        div            q,            vf00[w], \vertex[w]
        mul.xyz        \vertexoutput, \vertex, q
        move.w        \vertexoutput, \vertex
        mul            \stoutput,    \st,    q
        .endm

```

```

;///-----
;/// VertexFPtoGsXYZ2 - Convert an XYZW, floating-point vertex to GS
;/// XYZ2 format (ADC bit isnt set)
;///-----
    .macro    VertexFpToGsXYZ2    outputxyz,vertex
    ftoi4.xy        \outputxyz, \vertex
    ftoi0.z         \outputxyz, \vertex
    mfir.w          \outputxyz, vi00
    .endm

;///-----
;/// VertexFPtoGsXYZ2Adc - Convert an XYZW, floating-point vertex to GS
;/// XYZ2 format (ADC bit is set)
;///-----
    .macro    VertexFpToGsXYZ2Adc    outputxyz,vertex
    ftoi4.xy        \outputxyz, \vertex
    ftoi0.z         \outputxyz, \vertex
    ftoi15.w        \outputxyz, vf00
    .endm

;///-----
;/// VertexFpToGsXYZF2 - Convert an XYZF, floating-point vertex to GS
;/// XYZF2 format (ADC bit isnt set)
;///-----
    .macro    VertexFpToGsXYZF2    outputxyz,vertex
    ftoi4            \outputxyz, \vertex
    .endm

;///-----
;/// VertexFpToGsXYZF2Adc - Convert an XYZF, floating-point vertex to GS
;/// XYZF2 format (ADC bit is set)
;///-----
    .macro    VertexFpToGsXYZF2Adc    outputxyz,vertex
    ftoi4            \outputxyz, \vertex
    mtir            vclsmlitemp, \outputxyz[w]
    iaddiu          vclsmlitemp, 0x7FFF
    iaddi           vclsmlitemp, 1
    mfir.w          \outputxyz, vclsmlitemp
    .endm

;///-----
;/// ColorFPtoGsRGBAQ - Convert an RGBA, floating-point color to GS
;/// RGBAQ format
;///-----
    .macro    ColorFPtoGsRGBAQ    outputrgba,color
    ftoi0            \outputrgba, \color
    .endm

;///-----
;/// ColorGsRGBAQtoFP - Convert an RGBA, GS RGBAQ format to floating-
;/// point color
;///-----
    .macro    ColorGsRGBAQtoFP    outputrgba,color
    itof0            \outputrgba, \color
    .endm

```

```

;///-----
;/// CreateGsPRIM - Create a GS-packed-format PRIM command, according to
;/// a specified immediate value "prim"
;///
;/// Note: Meant more for debugging purposes than for a final solution
;///-----
        .macro    CreateGsPRIM    outputprim,prim
        iaddiu    vclsmllitemp, vi00, \prim
        mfir      \outputprim, vclsmllitemp
        .endm

;///-----
;/// CreateGsRGBA - Create a GS-packed-format RGBA command, according to
;/// specified immediate values "r", "g", "b" and "a" (integer 0-255)
;///
;/// Note: Meant more for debugging purposes than for a final solution
;///-----
        .macro    CreateGsRGBA    outputrgba,r,g,b,a
        iaddiu    vclsmllitemp, vi00, \r
        mfir.x    \outputrgba, vclsmllitemp
        iaddiu    vclsmllitemp, vi00, \g
        mfir.y    \outputrgba, vclsmllitemp
        iaddiu    vclsmllitemp, vi00, \b
        mfir.z    \outputrgba, vclsmllitemp
        iaddiu    vclsmllitemp, vi00, \a
        mfir.w    \outputrgba, vclsmllitemp
        .endm

;///-----
;/// CreateGsSTQ - Create a GS-packed-format STQ command, according to
;/// specified immediate values "s", "t" and "q" (floats)
;///
;/// Note: I register is modified
;/// Note: Meant more for debugging purposes than for a final solution
;///-----
        .macro    CreateGsSTQ    outputstq,s,t,q
        loi       \s
        add.x     \outputstq, vf00, i
        loi       \t
        add.y     \outputstq, vf00, i
        loi       \q
        add.z     \outputstq, vf00, i
        .endm

;///-----
;/// CreateGsUV - Create a GS-packed-format VU command, according to
;/// specified immediate values "u" and "v" (integer -32768 - 32768,
;/// with 4 LSB as precision)
;///
;/// Note: Meant more for debugging purposes than for a final solution
;///-----
        .macro    CreateGsUV    outputuv,u,v
        iaddiu    vclsmllitemp, vi00, \u
        mfir.x    \outputuv, vclsmllitemp
        iaddiu    vclsmllitemp, vi00, \v
        mfir.y    \outputuv, vclsmllitemp
        .endm

```

```

;///-----
;/// CreateGsRGBA - Create a GS-packed-format RGBA command, according to
;/// a specified immediate value "fog" (integer 0-255)
;///
;/// Note: Meant more for debugging purposes than for a final solution
;///-----
        .macro    CreateGsFOG    outputfog,fog
        iaddiu    vclsmlltemp, vi00, \fog * 16
        mfiw     \outputfog, vclsmlltemp
        .endm

;///-----
;/// VectorDotProduct - Calculate the dot product of "vector1" and
;/// "vector2", and output to "dotproduct"[x]
;///-----
        .macro    VectorDotProduct    dotproduct,vector1,vector2
        mul.xyz   \dotproduct, \vector1, \vector2
        add.x     \dotproduct, \dotproduct, \dotproduct[y]
        add.x     \dotproduct, \dotproduct, \dotproduct[z]
        .endm

;///-----
;/// VectorDotProductACC - Calculate the dot product of "vector1" and
;/// "vector2", and output to "dotproduct"[x]. This one does it using
;/// the ACC register which, depending on the case, might turn out to be
;/// faster or slower.
;///
;/// Note: ACC register is modified
;///-----
        .macro    VectorDotProductACC    dotproduct,vector1,vector2
        max       Vector1111, vf00, vf00[w]
        mul       vclsmllftemp, \vector1, \vector2
        add.x     acc, vclsmllftemp, vclsmllftemp[y]
        madd.x    \dotproduct, Vector1111, vclsmllftemp
        .endm

;///-----
;/// VectorCrossProduct - Calculate the cross product of "vector1" and
;/// "vector2", and output to "vectoroutput"
;///
;/// Note: ACC register is modified
;///-----
        .macro    VectorCrossProduct    vectoroutput,vector1,vector2
        opmula.xyz    ACC, \vector1, \vector2
        opmsub.xyz    \vectoroutput, \vector2, \vector1
        sub.w         \vectoroutput, vf00, vf00
        .endm

```

```

;///-----
;/// VectorNormalize - Bring the length of "vector" to 1.f, and output
;/// it to "vectoroutput"
;///
;/// Note: Q register is modified
;///-----
    .macro    VectorNormalize    vecoutput,vector
    mul.xyz    vclsmfltemp, \vector,    \vector
    add.x      vclsmfltemp, vclsmfltemp, vclsmfltemp[y]
    add.x      vclsmfltemp, vclsmfltemp, vclsmfltemp[z]
    rsqrt     q,                vf00[w],    vclsmfltemp[x]
    sub.w     \vecoutput,    vf00,        vf00
    mul.xyz    \vecoutput,    \vector,    q
    .endm

;///-----
;/// VectorNormalizeXYZ - Bring the length of "vector" to 1.f, and out-
;/// put it to "vectoroutput". The "w" field isn't transferred.
;///
;/// Note: Q register is modified
;///-----
    .macro    VectorNormalizeXYZ    vecoutput,vector
    mul.xyz    vclsmfltemp, \vector,    \vector
    add.x      vclsmfltemp, vclsmfltemp, vclsmfltemp[y]
    add.x      vclsmfltemp, vclsmfltemp, vclsmfltemp[z]
    rsqrt     q,                vf00[w],    vclsmfltemp[x]
    mul.xyz    \vecoutput,    \vector,    q
    .endm

;///-----
;/// VertexLightAmb - Apply ambient lighting "ambientrgba" to a vertex
;/// of color "vertexrgba", and output the result in "outputrgba"
;///-----
    .macro    VertexLightAmb    rgbaout,vertexrgba,ambientrgba
    mul        \rgbaout, \vertexrgba, \ambientrgba
    .endm

;///-----
;/// VertexLightDir3 - Apply up to 3 directional lights contained in a
;/// light matrix "lightmatrix" to a vertex of color "vertexrgba" and
;/// having a normal "vertexnormal", and output the result in
;/// "outputrgba"
;///
;/// Note: ACC register is modified
;///-----
    .macro    VertexLightDir3
    rgbaout,vertexrgba,vertexnormal,lightcolors,lightnormals
    mul        acc,        \lightnormals[0], \vertexnormal[x]
    madd       acc,        \lightnormals[1], \vertexnormal[y]
    madd       acc,        \lightnormals[2], \vertexnormal[z]
    madd       \rgbaout, \lightnormals[3], \vertexnormal[w] ;// Here
    "rgbaout" is the dot product for the 3 lights
    max        \rgbaout, \rgbaout,        vf00[x] ;// Here
    "rgbaout" is the dot product for the 3 lights
    mul        acc,        \lightcolors[0], \rgbaout[x]
    madd       acc,        \lightcolors[1], \rgbaout[y]
    madd       \rgbaout, \lightcolors[2], \rgbaout[z] ;// Here
    "rgbaout" is the light applied on the vertex
    mul        \rgbaout, \vertexrgba,    \rgbaout ;// Here
    "rgbaout" is the amount of light reflected by the vertex
    .endm

```

```

;///-----
;/// VertexLightDir3Amb - Apply up to 3 directional lights, plus an
;/// ambient light contained in a light matrix "lightmatrix" to a vertex
;/// of color "vertexrgba" and having a normal "vertexnormal", and
;/// output the result in "outputrgba"
;///
;/// Note: ACC register is modified
;///-----
        .macro    VertexLightDir3Amb
rgbaout,vertexrgba,vertexnormal,lightcolors,lightnormals
        mul        acc,        \lightnormals[0], \vertexnormal[x]
        madd       acc,        \lightnormals[1], \vertexnormal[y]
        madd       acc,        \lightnormals[2], \vertexnormal[z]
        madd       \rgbaout, \lightnormals[3], \vertexnormal[w] ;// Here
"rgbaout" is the dot product for the 3 lights
        max        \rgbaout, \rgbaout,          vf00[x]          ;// Here
"rgbaout" is the dot product for the 3 lights
        mul        acc,        \lightcolors[0], \rgbaout[x]
        madd       acc,        \lightcolors[1], \rgbaout[y]
        madd       acc,        \lightcolors[2], \rgbaout[z]
        madd       \rgbaout, \lightcolors[3], \rgbaout[w]          ;// Here
"rgbaout" is the light applied on the vertex
        mul.xyz    \rgbaout, \vertexrgba,      \rgbaout          ;// Here
"rgbaout" is the amount of light reflected by the vertex
        .endm

;///-----
;/// FogSetup - Set up fog "fogparams", by specifying "nearfog" and
;/// "farfog". "fogparams" will afterward be ready to be used by fog-
;/// related macros, like "VertexFogLinear" for example.
;///
;/// Note: I register is modified
;///-----
        .macro    FogSetup fogparams,nearfogz,farfogz
        sub        \fogparams, vf00,          vf00          ;// Set XYZW to
0
        loi        \farfogz
        add.w      \fogparams, \fogparams,    i              ;// fogparam[w]
is farfogz
        loi        \nearfogz
        add.z      \fogparams, \fogparams,    \fogparams[w]
        sub.z      \fogparams, \fogparams,    i
        loi        255.0
        add.xy     \fogparams, \fogparams,    i              ;// fogparam[y]
is 255.0
        sub.x      \fogparams, \fogparams,    vf00[w]        ;// fogparam[x]
is 254.0
        div        q,        \fogparams[y], \fogparams[z]
        sub.z      \fogparams, \fogparams,    \fogparams
        add.z      \fogparams, \fogparams,    q;// fogparam[z] is 255.f /
(farfogz - nearfogz)
        .endm

```

```

;///-----
;/// VertexFogLinear - Apply fog "fogparams" to a vertex "xyzw", and
;/// output the result in "xyzfoutput". "xyzw" [w] is assumed to be
;/// the distance from the camera. "fogparams" must contain farfogz in
;/// [w], and (255.f / (farfogz - nearfogz)) in [z]. "xyzfoutputf" [w]
;/// will contain a float value between 0.0 and 255.0, inclusively.
;///-----
    .macro    VertexFogLinear    xyzfoutput,xyzw,fogparams
    move.xyz    \xyzfoutput, \xyzw                ;// XYZ part won't
be modified
    sub.w      \xyzfoutput, \fogparams, \xyzw[w]    ;// fog = (farfogz -
z) * 255.0 /
    mul.w      \xyzfoutput, \xyzfoutput, \fogparams[z];//    (farfogz -
nearfogz)
    max.w      \xyzfoutput, \xyzfoutput, vf00[x]    ;// Clamp fog values
outside the
    mini.w     \xyzfoutput, \xyzfoutput, \fogparams[y];// range 0.0-255.0
    .endm

;///-----
;/// VertexFogRemove - Remove any effect of fog to "xyzf". "fogparams"
;/// [x] must be set to 254.0. "xyzf" will be modified directly.
;///-----
    .macro    VertexFogRemove    xyzf,fogparams
    add.w      \xyzf, vf00, \fogparams[x]    ;// xyzw[w] = 1.0 + 254.0 =
255.0 = no fog
    .endm

;///-----
;/// PushInteger1 - Push "integer1" on "stackptr"
;///
;/// Note: "stackptr" is updated
;///-----
    .macro    PushInteger1    stackptr,integer1
    isubiu     \stackptr, \stackptr, 1
    iswr.x     \integer1, (\stackptr):VCLSML_STACK
    .endm

;///-----
;/// PushInteger2 - Push "integer1" and "integer2" on "stackptr"
;///
;/// Note: "stackptr" is updated
;///-----
    .macro    PushInteger2    stackptr,integer1,integer2
    isubiu     \stackptr, \stackptr, 1
    iswr.x     \integer1, (\stackptr):VCLSML_STACK
    iswr.y     \integer2, (\stackptr):VCLSML_STACK
    .endm

;///-----
;/// PushInteger3 - Push "integer1", "integer2" and "integer3" on
;/// "stackptr"
;///
;/// Note: "stackptr" is updated
;///-----
    .macro    PushInteger3    stackptr,integer1,integer2,integer3
    isubiu     \stackptr, \stackptr, 1
    iswr.x     \integer1, (\stackptr):VCLSML_STACK
    iswr.y     \integer2, (\stackptr):VCLSML_STACK
    iswr.z     \integer3, (\stackptr):VCLSML_STACK
    .endm

```

```

;///-----
;/// PushInteger4 - Push "integer1", "integer2", "integer3" and
;/// "integer4" on "stackptr"
;///
;/// Note: "stackptr" is updated
;///-----
    .macro    PushInteger4    stackptr,integer1,integer2,integer3,integer4
    isubiu    \stackptr, \stackptr, 1
    iswr.x    \integer1, (\stackptr):VCLSML_STACK
    iswr.y    \integer2, (\stackptr):VCLSML_STACK
    iswr.z    \integer3, (\stackptr):VCLSML_STACK
    iswr.w    \integer4, (\stackptr):VCLSML_STACK
    .endm

;///-----
;/// PopInteger1 - Pop "integer1" on "stackptr"
;///
;/// Note: "stackptr" is updated
;///-----
    .macro    PopInteger1    stackptr,integer1
    ilwr.x    \integer1, (\stackptr):VCLSML_STACK
    iaddiu    \stackptr, \stackptr, 1
    .endm

;///-----
;/// PopInteger2 - Pop "integer1" and "integer2" on "stackptr"
;///
;/// Note: "stackptr" is updated
;///-----
    .macro    PopInteger2    stackptr,integer1,integer2
    ilwr.y    \integer2, (\stackptr):VCLSML_STACK
    ilwr.x    \integer1, (\stackptr):VCLSML_STACK
    iaddiu    \stackptr, \stackptr, 1
    .endm

;///-----
;/// PopInteger3 - Pop "integer1", "integer2" and "integer3" on
;/// "stackptr"
;///
;/// Note: "stackptr" is updated
;///-----
    .macro    PopInteger3    stackptr,integer1,integer2,integer3
    ilwr.z    \integer3, (\stackptr):VCLSML_STACK
    ilwr.y    \integer2, (\stackptr):VCLSML_STACK
    ilwr.x    \integer1, (\stackptr):VCLSML_STACK
    iaddiu    \stackptr, \stackptr, 1
    .endm

```

```

;///-----
;/// PopInteger4 - Pop "integer1", "integer2", "integer3" and
;/// "integer4" on "stackptr"
;///
;/// Note: "stackptr" is updated
;///-----
    .macro    PopInteger4    stackptr,integer1,integer2,integer3,integer4
ilwr.w        \integer4, (\stackptr):VCLSML_STACK
ilwr.z        \integer3, (\stackptr):VCLSML_STACK
ilwr.y        \integer2, (\stackptr):VCLSML_STACK
ilwr.x        \integer1, (\stackptr):VCLSML_STACK
iaddiu        \stackptr, \stackptr, 1
    .endm

;///-----
;/// PushMatrix - Push "matrix" onto the "stackptr"
;///
;/// Note: "stackptr" is updated
;///-----
    .macro    PushMatrix    stackptr,matrix
sq            \matrix[0], -1(\stackptr):VCLSML_STACK
sq            \matrix[1], -2(\stackptr):VCLSML_STACK
sq            \matrix[2], -3(\stackptr):VCLSML_STACK
sq            \matrix[3], -4(\stackptr):VCLSML_STACK
iaddi        \stackptr, \stackptr, -4
    .endm

;///-----
;/// PopMatrix - Pop "matrix" out of the "stackptr"
;///
;/// Note: "stackptr" is updated
;///-----
    .macro    PopMatrix     stackptr,matrix
lq            \matrix[0], 0(\stackptr):VCLSML_STACK
lq            \matrix[1], 1(\stackptr):VCLSML_STACK
lq            \matrix[2], 2(\stackptr):VCLSML_STACK
lq            \matrix[3], 3(\stackptr):VCLSML_STACK
iaddi        \stackptr, \stackptr, 4
    .endm

;///-----
;/// PushVector - Push "vector" onto the "stackptr"
;///
;/// Note: "stackptr" is updated
;///-----
    .macro    PushVector    stackptr,vector
sqd          \vector, (--\stackptr):VCLSML_STACK
    .endm

;///-----
;/// PopVector - Pop "vector" out of the "stackptr"
;///
;/// Note: "stackptr" is updated
;///-----
    .macro    PopVector     stackptr,vector
lqi          \vector, (\stackptr++):VCLSML_STACK
    .endm

```

```

;///-----
;/// PushVertex - Push "vector" onto the "stackptr"
;///
;/// Note: "stackptr" is updated
;///-----
    .macro    PushVertex    stackptr,vertex
    sqd      \vertex, (--\stackptr):VCLSML_STACK
    .endm

;///-----
;/// PopVertex - Pop "vertex" out of the "stackptr"
;///
;/// Note: "stackptr" is updated
;///-----
    .macro    PopVertex    stackptr,vertex
    lqi      \vertex, (\stackptr++):VCLSML_STACK
    .endm

;///-----
;/// AngleSinCos - Returns the sin and cos of up to 2 angles, which must
;/// be contained in the X and Z elements of "angle". The sin/cos pair
;/// will be contained in the X/Y elements of "sincos" for the first
;/// angle, and Z/W for the second one.
;/// Thanks to Colin Hugues (SCEE) for that one
;///
;/// Note: ACC and I registers are modified, and a bunch of temporary
;/// variables are created... Maybe bad for VCL register pressure
;///-----
    .macro    AngleSinCos    angle,sincos
    move.xz  \sincos, \angle                ; To avoid modifying the
original angles...

    mul.w    \sincos, vf00, \sincos[z]      ; Copy angle from z to w
    add.y    \sincos, vf00, \sincos[x]      ; Copy angle from x to y

    loi     1.570796                        ; Phase difference for sin as
cos ( PI/2 )
    sub.xz  \sincos, \sincos, I             ;

    abs     \sincos, \sincos                ; Mirror cos
around zero

    max     Vector1111, vf00, vf00[w]       ; Initialise all
1s

    loi     -0.159155                        ; Scale so single cycle is range 0 to -1 (
*-1/2PI )
    mul     ACC, \sincos, I                 ;

    loi     12582912.0                       ; Apply bias to remove
fractional part
    msub    ACC, Vector1111, I               ;
    madd    ACC, Vector1111, I               ; Remove bias to leave
original int part

    loi     -0.159155                        ; Apply original number to leave
fraction range only
    msub    ACC, \sincos, I                 ;

    loi     0.5                              ; Ajust range: -
0.5 to +0.5

```

```

msub          \sincos, Vector1111, I          ;
abs           \sincos, \sincos                ; Clamp: 0 to
+0.5
loi          0.25                             ; Adjust range: -
0.25 to +0.25
sub          \sincos, \sincos, I            ;
mul          anglepower2, \sincos, \sincos    ; a^2
loi          -76.574959                       ;
mul          k4angle, \sincos, I             ; k4 a
loi          -41.341675                       ;
mul          k2angle, \sincos, I             ; k2 a
loi          81.602226                       ;
mul          k3angle, \sincos, I             ; k3 a
mul          anglepower4, anglepower2, anglepower2 ; a^4
mul          k4angle, k4angle, anglepower2    ; k4 a^3
mul          ACC, k2angle, anglepower2       ; + k2 a^3
loi          39.710659                       ; k5 a
mul          k2angle, \sincos, I             ;
mul          anglepower8, anglepower4, anglepower4 ; a^8
madd         ACC, k4angle, anglepower4       ; + k4 a^7
madd         ACC, k3angle, anglepower4       ; + k3 a^5
loi          6.283185                         ;
madd         ACC, \sincos, I                 ; + k1 a
madd         \sincos, k2angle, anglepower8   ; + k5 a^9
.endm

;/-------
; // QuaternionToMatrix - Converts a quaternion rotation to a matrix
; // Thanks to Colin Hugues (SCEE) for that one
; //
; // Note: ACC and I registers are modified
; //-----
.macro QuaternionToMatrix matresult,quaternion

mula.xyz     ACC, \quaternion, \quaternion    ; xx yy zz

loi          1.414213562
muli         vclsm1ftmp, \quaternion, I      ; x sqrt2 y sqrt2 z
sqrt2 w sqrt2

mr32.w      \matresult[0], vf00              ; Set rhs matrix
line 0 to 0
mr32.w      \matresult[1], vf00              ;
mr32.w      \matresult[2], vf00              ; Set rhs matrix
move        \matresult[3], vf00              ; Set bottom line to
0 0 0 1

madd.xyz    vcl_2qq, \quaternion, \quaternion ; 2xx      2yy
2zz
addw.xyz    Vector111, vf00, vf00            ; 1        1
1          -

```

```

    opmula.xyz    ACC,  vclsmftemp, vclsmftemp      ; 2yz      2xz
2xy      -
    msubw.xyz    vclsmftemp2, vclsmftemp, vclsmftemp; 2yz-2xw  2xz-2yz
2xy-2zw   -
    maddw.xyz    vclsmftemp3, vclsmftemp, vclsmftemp; 2yz+2xw  2xz+2yz
2xy+2zw   -
    addaw.xyz    ACC,  vf00, vf00                  ; 1          1
1          -
    msubax.xyz   ACC,  Vector111, vcl_2qq          ; 1          1-2xx
1-2xx     -

    msuby.z      \matresult[2], Vector111, vcl_2qq  ; -          -
1-2xx-2yy -
    msubay.x     ACC,  Vector111, vcl_2qq          ; 1-2yy      1-2xx
1-2xx-2yy -
    msubz.y      \matresult[1], Vector111, vcl_2qq  ; -          1-2xx-
2zz -
-
    mr32.y       \matresult[0], vclsmftemp2
    msubz.x      \matresult[0], Vector111, vcl_2qq  ; 1-2yy-2zz -
-
    mr32.x       \matresult[2], vclsmftemp2
    addy.z       \matresult[0], vf00, vclsmftemp3
    mr32.w       vclsmftemp, vclsmftemp2
    mr32.z       \matresult[1], vclsmftemp
    addx.y       \matresult[2], vf00, vclsmftemp3
    mr32.y       vclsmftemp3, vclsmftemp3
    mr32.x       \matresult[1], vclsmftemp3

    .endm

;///-----
;/// QuaternionMultiply - Multiplies "quaternion1" and "quaternion2",
;/// and puts the result in "quatresult".
;/// Thanks to Colin Hugues (SCEE) for that one
;///
;/// Note: ACC register is modified
;///-----
    .macro   QuaternionMultiply  quatresult,quaternion1,quaternion2
    mul      vclsmftemp, \quaternion1, \quaternion2  ; xx yy zz ww

    Outerproduct
    opmula.xyz    ACC,          \quaternion1, \quaternion2  ; Start
    madd.xyz      ACC,          \quaternion1, \quaternion2[w]; Add w2.xyz1
    madd.xyz      ACC,          \quaternion2, \quaternion1[w]; Add w1.xyz2
    opmsub.xyz    \quatresult, \quaternion2, \quaternion1  ; Finish
    Outerproduct

    sub.w         ACC,          vclsmftemp,  vclsmftemp[z] ; ww - zz
    msub.w        ACC,          vf00,        vclsmftemp[y] ; ww - zz - yy
    msub.w        \quatresult, vf00,        vclsmftemp[x] ; ww - zz - yy -
xx
    .endm

```

## Appendix B: Detailed Information Regarding Loops in VCL

### Pipelining and VCL

When the “`--LoopCS`” directive is used, VCL will attempt to unroll and pipeline back to the start any block of code that ends with a conditional branch.

VCL analyzes the sequence of instructions and determines what the best size for the loop would be. A simplistic calculation for the loop size is:

```
best_size = max (number_upper_instructions, number_lower_instructions,
sum_throughput_p, sum_throughput_q)
```

The actual calculation is actually more involved, particularly due to issues like IALU instructions placing (with respect to the branch) and the possibility of circular dependency chains in the instruction sequence.

**Note:** VCL currently relies on the user to identify memory store/load sequences. It does so by the use of tags. (Refer to “Memory Aliasing and Instructions Reordering” for more details.)

For a typical renderer, the usual circular chain is of the form:

```
IADDIU      ptr, ptr, sizeof
IBNE        ptr, end_ptr, loop
```

Or:

```
IADDI       count, count, -1
IBNE        count, vi00, loop
```

These cases aren’t long compared to the size of the loop. (The rest of a typical renderer pulls data out of memory, processes it, then writes it back to memory. However, it does not cross a loop iteration boundary.)

For other types of code loops (such as physics dynamics calculations on strings), where the output from one loop iteration is an input to the next, and the length of the calculation is the same as the overall length of the loop, pipelining will not greatly improve performances without a modification of the algorithm. Such a modification could be to process multiple strings at once.

After analyzing the code, VCL tries to schedule it so that it fits into a block that is of size greater or equal to 0, but smaller than the loop size. It does so by wrapping instructions off the end and back to the start. The number of times an instruction is wrapped around the loop will determine the stage number to which it belongs.

In short:

- In linear mode (non-looped), VCL schedules in  $Z$  instructions.
- In loop mode, it schedules in  $Z$  modulo  $n$ , where  $n$  equals `best_size + current optimization phase`.
- Software pipelining where the branch is at the end of the critical path will not be optimized greatly by VCL.

Study the following example, which will be referred to later in this appendix:

```
vrt_loop:
  --LoopCS      10,10
  LQI           vrt, (in_p++)           ; Load vertex

  MUL           acc, mat[0], vrt[x]     ; Transform vertex
  MADD          acc, mat[1], vrt[y]     ;
```

```

MADD      acc, mat[2], vrt[z]      ;
MADD      camv, mat[3], vf00[w]   ;

DIV       q, vf00[w], camv[w]     ; Perspective correction
MUL      screenv, camv, q         ;

FTOI4    fixpt, screenv           ; Convert to GS format

SQI      fixpt, (out_p++)         ; Save the vertex

IADDI    count, count, -1         ; Next vertex...
IBNE     count, vi00, vrt_loop   ;
    
```

The following is a typical sequence of instructions for a loop, where the numerals 1 to 9 denote blocks of instructions that are grouped by pipeline stages. Execution in the order 1, 2, 3, 4, 5, 6, 7, 8, 9 is equivalent to one iteration of the loop.

Figure 1

1
2
3
4
5
6c
7
8
9m

“c” denotes the stage for the conditional.. In the example above, it is highly likely that the conditional will be at stage 1, since it is not dependent on much else in the code. “m” is the last stage. In a typical renderer, this would often correspond to the color calculations.

As pipeline execute happens:

Figure 2

Prologue	1									P1 starts
	2	1								P2 starts, P1 is at stage 2
	3	2	1							P3 starts, P2 is at stage 2, P3 is at stage 3
	4	3	2	1						Etc...
	5	4	3	2	1					
	6c	5	4	3	2	1				
	7	6c	5	4	3	2	1			
	8	7	6c	5	4	3	2	1		
MainLoop	9	8	7	6c	5	4	3	2	1	Main body of the loop
Epilogue		9	8	7						
			9	8						
				9						

If the conditional is at stage C (here, 6), then stages 1 to C-1 will miss-execute, i.e. in the main loop, the pipeline runs C-1 stages ahead of the conditional. But when the condition is found to be true, VCL will only complete the processing for valid stages which are in the graph above 9, 8-9, and 7-8-9 (epilogue part).

## “--LoopCS n,m” Directive

### n (Minimum Number of Loops)

For small loop counts, there are many instructions that are associated with pipeline stages ahead of the conditional. To get better performance for a small count ( $n$  greater or equal to 1 but smaller than  $M-C$ ), once the conditional in the loop is encountered, VCL can jump to a special case to complete the calculations on the required pipeline stages for this count.

Following is a modified version of the above diagram:

Figure 3

	1								
	2	1							
	3	2	1						
	4	3	2	1					
	5	4	3	2	1				
A	6c	5	4	3	2	1			
B	7	6c	5	4	3	2	1		
C	8	7	6c	5	4	3	2	1	
MainLoop	9	8	7	6c	5	4	3	2	1
		9	8	7					
			9	8					
				9					

If the code has a small minimum count, such as  $n = 1$ , it is possible for the code to exit at “A” and go to a special-case epilogue (EPI\_A). The steps that have already taken place are:

Figure 4

	1						
	2	1					
	3	2	1				
	4	3	2	1			
	5	4	3	2	1		
A	6c	5	4	3	2	1	

If the condition at stage 6 is found to be true, the following will be executed:

Figure 5

EPI_A	7						
	8						
	9						

Similarly for  $n = 2$ , a special epilogue may be created for the following (EPI-B):

Figure 6

EPI_B	8	7					
	9	8					
		9					

In the case above, having nine stages would require the creation of nine special epilogues, which is a lot of code generation. However, if VCL is told –via the “Minimum Number of Loops”- that there will always be, for example, three iterations, then the special case codes EPI\_A and EPI\_B as well as the two conditionals A and B may be removed altogether.

Up to at least version 1.3, VCL doesn’t reschedule instructions across conditionals. So conditional removals by ways described above will most certainly result in better code optimization.

**m (Slop Count)**

Referring to the tables above, it can be seen that, in the main loop, VCL will execute some stages ahead of the conditional (in the example, stages 1 to 5). If these contain instructions with side effects (like memory stores and XGKick), this could result in data corruption, since by the time the conditional takes place, such instructions would have already executed.

In the following table, “\*” denotes stages containing an instruction with side-effects, and “f” denotes the first instruction containing instructions with side-effects. “c” denotes the conditional stage, and “m” the last stage.

**Figure 7**

1
2f*
3
4*
5
6c
7*
8
9m

If the first side effect is at stage 2, then stage 2 will have miss-executed a maximum count greater or equal to 0, but smaller than c-f (6-2=4). This is acceptable if extra padding is provided at the end of the store buffer. The number of available padding slots is specified with the Slop Count.

If, for example, m=1, then the above case would generate incorrect code, since c-f=4.

However, for the following case, the generated code would be correct, since c-f=1.

**Figure 8**

1
2
3
4
5f*
6c
7*
8
9m

If no side-effect stages can be mis-executed, then the Slop Count must be set to 0. Note that this will, however, result in real constraints on VCL code generation.

**.rawloop / .endrawloop**

For regular loops, VCL will analyze the loop and decide how to stage the instructions and reschedule them accordingly. Then compatible prologues and epilogues will be created around the loop's main body.

In some cases, however, you may already know what the stages are like, and simply want VCL to unroll them and create the prologue and epilogue. Use raw loops for this..

The regular loop shown in "Pipelining and VCL" would be similar to the following example, using raw loops:

```
.rawloop

vrt_loop:
    --LoopCS      10,10

    2..MADD      acc, mat[2], vrt[z]    ; Rotate vertex .. 2
= 1..LQI       vrt, (in_p++)          ; Load vertex, increment pointer

    2..MADD      camv, mat[3], vf00[w] ; Translate rotated vertex
= 3..MOVE      ocamv, camv           ; Save copy of camera space coordinate

    NOP
= 5..SQI       fixpt, (out_p++)       ; Save out GS-format vertex

    4..FTOI4     fixpt, screenv        ; Convert screen coordinate to GS-format
= 1..IADDI     count, count, -1       ; Decrement loop counter

    1..MUL       acc, mat[0], vrt[x]   ; Rotate vertex .. 0
= NOP

    1..MADD      acc, mat[0], vrt[y]   ; Rotate vertex .. 1
= 1..IBNE     count, vi00, vrt_loop   ; Reached the end?

    3..MUL       screenv, ocamv, q     ; Do perspective divide
    2..DIV       q, vf00[w], camv[w]   ; Start perspective divide calculation

.endrawloop
```

The "<n>.." instruction syntax tells VCL in which stage of the loop the instruction belongs, so it can generate proper prologues and epilogues.

In raw mode, "=" may be used as a "line continue" character, as long as it is the first non-white character on a lower-instruction line. This permits better comments placement. If not used, upper and lower instructions must be placed on the same line, much like regular VSM code.

VCL will create the prologue by first taking the instructions with a "1.." suffix, then the instructions with a "2.." suffix, and so on.

Prologue and epilogue instructions are not rescheduled in the case of a raw loop, as they are for regular loop unrolling. Therefore, this may result in sub-optimal code. However, this is necessary for cases where the code needs to run on VU0 in parallel with code on the EE, without synchronization.

